

Random Euclidean Bipartite Matching Problems

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- 3 Concentration inequalities
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Setting of the problem

Let $d \geq 1$, and consider $2n \geq 1$ i.i.d. random variables

$$(X_i)_{i=1}^n, \quad (Y_i)_{i=1}^n$$

taking values in $[0, 1]^d$ with **uniform distribution**, i.e.,

$$P(X_i \in A) = |A| = \int_A dx, \quad \text{for every } A \subseteq [0, 1]^d \text{ Borel,}$$

The **Random Euclidean Bipartite Matching Problem** is defined as the following random variational problem,

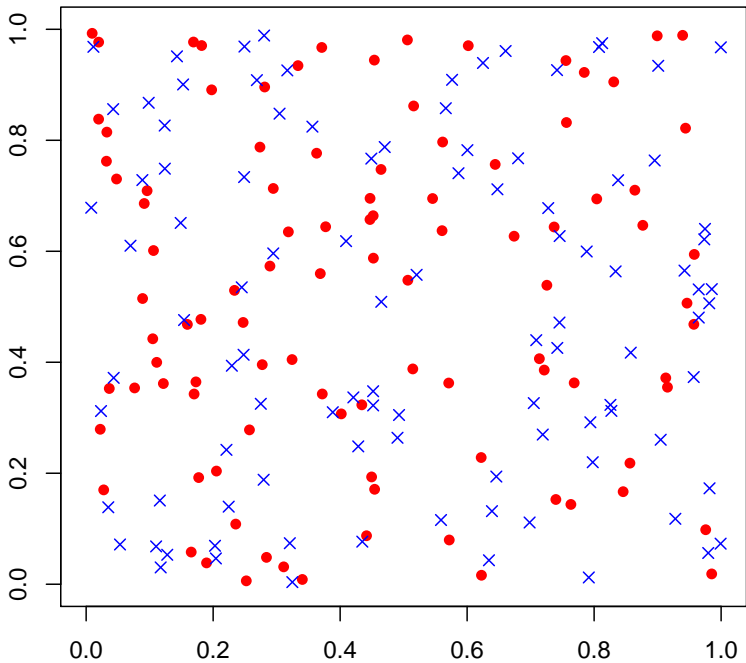
$$\min_{\sigma \in \mathcal{S}^n} \sum_{i=1}^n |X_i - Y_{\sigma(i)}|,$$

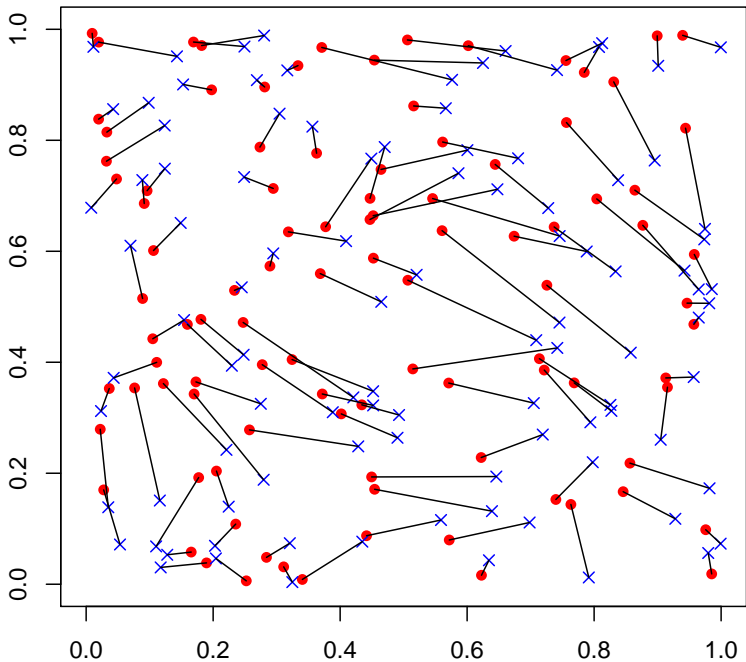
where \mathcal{S}^n denotes the set of permutations over $\{1, \dots, n\}$.

Variant: power of distance

$$\text{For } p \in (0, +\infty), \quad B_{p,n} := \min_{\sigma \in \mathcal{S}^n} \sum_{i=1}^n |X_i - Y_{\sigma(i)}|^p.$$

Let us look at a simulation with $n = 100$, $d = 2$, $p = 2$.





We can devise lots of variants:

- 1 Different domains
- 2 Different laws
- 3 Different distances (costs)
- 4 Different number of red/blue points: what is a matching?

Another related problem is the **(monopartite) minimal matching problem**:

Let $(X_i)_{i=1}^{2n}$ be i.i.d. uniform and define

$$M_{p,n} = \min_{\sigma \in \mathcal{S}^{2n}} \sum_{i=1}^n |X_{\sigma(i)} - X_{\sigma(n+i)}|^p,$$

that is, we forget about colouring of the points.

Exercise

Show that $\mathbb{E}[M_{p,n}] \leq \mathbb{E}[B_{p,n}]$.

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Heuristics

Points are “uniformly distributed” on $[0, 1]^d$
 \Rightarrow for each X_i we find Y_j at a distance $\approx n^{-1/d}$.

Considering the p -th power and summing upon n leads to

$$B_{p,n} \approx n \cdot \frac{1}{n^{p/d}} = n^{1-p/d}.$$

Conjecture

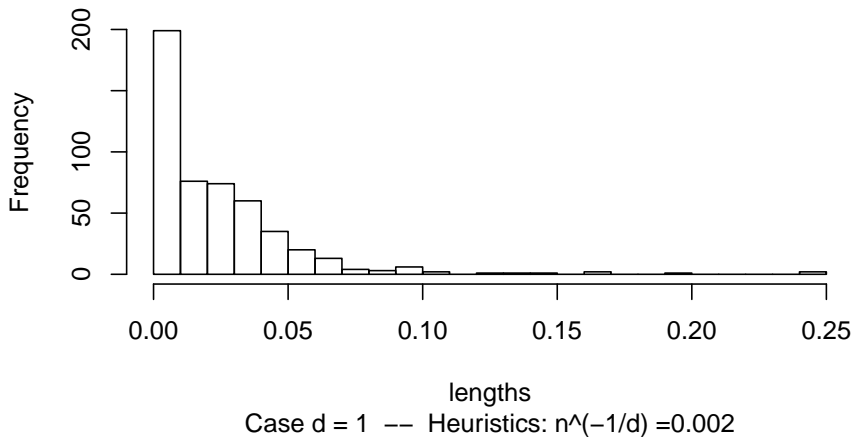
The limit (e.g. in probability)

$$\lim_{n \rightarrow \infty} \frac{B_{p,n}}{n^{1-p/d}}$$

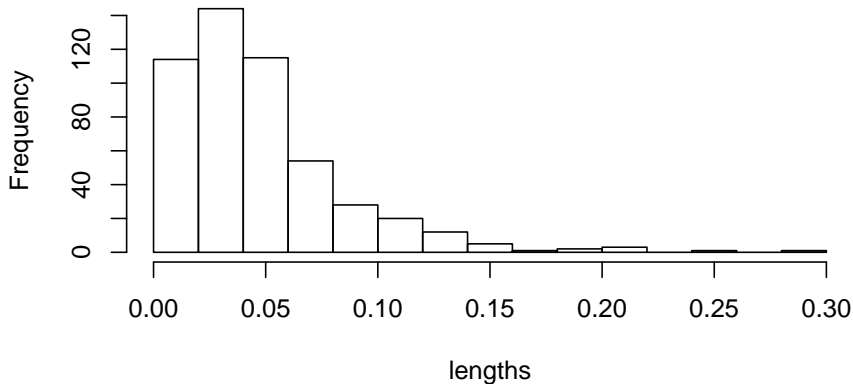
exists finite and strictly positive.

Histograms of (simulated) matching lengths with $p = 1$, $n = 500$ and $d \in \{1, 2, 3\}$.

Simulations: $d = 1, p = 1, n = 500$

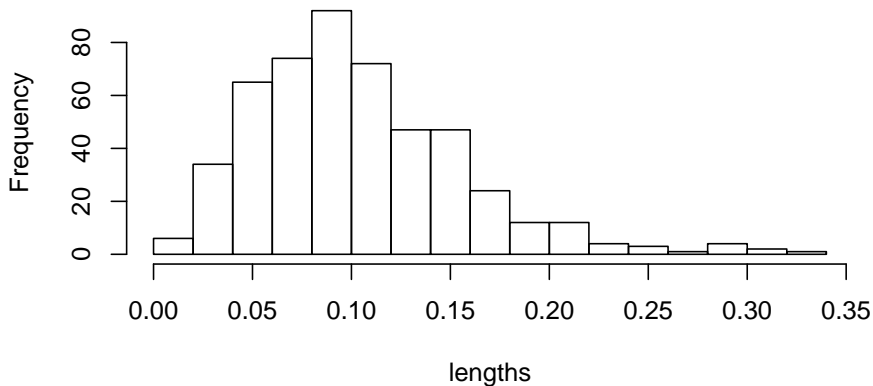


Simulations: $d = 2, p = 1, n = 500$



Case $d = 2$ -- Heuristics: $n^{(-1/d)} = 0.0447213595499958$

Simulations: $d = 3, p = 1, n = 500$



Case $d = 3$ --- Heuristics: $n^{-(1/d)} = 0.125992104989487$

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Proposition

There exists $c(d, p) > 0$ such that for every $n \geq 1$,

$$\mathbb{E}[B_{p,n}] \geq c(d, p)n^{1-p/d}.$$

We use the following lemma.

Lemma

Let $m \geq 1$ and $(X_i)_{i=1}^m$ i.i.d. on $[0, 1]^d$ with uniform distribution. Then, for every $x \in [0, 1]^d$,

$$\mathbb{E} \left[\min_{i=1, \dots, m} |X_i - x|^p \right] \geq \frac{c(d, p)}{m^{p/d}},$$

where $c(d, p) > 0$ depends on d and p only.

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where $c(d, p) > 0$ is depends on d and p only.

Proof of Lemma

Consider the **survival** (or reliability) function, i.e., for $t \in [0, \infty)$,

$$\begin{aligned} P\left(\min_{i=1, \dots, m} |X_i - x|^p > t\right) &= P(|X_1 - x| > t^{1/p})^m \quad \text{by independence} \\ &= \mathcal{L}^d\left([0, 1]^d \setminus \overline{B(x, t^{1/p})}\right)^m \quad \text{since the law is uniform} \\ &\geq \left(\max\left\{(1 - \omega_d t^{d/p}), 0\right\}\right)^m \end{aligned}$$

with $\omega_d = |B(0, 1)|$.

The general formula (for non-negative random variables Z)

$$\mathbb{E}[Z] = \int_0^\infty P(Z > t) dt$$

gives

$$\begin{aligned} \mathbb{E}\left[\min_{i=1, \dots, m} |X_i - x|^p\right] &= \int_0^\infty P\left(\min_{i=1, \dots, m} |X_i - x|^p > t\right) dt \\ &\geq \int_0^{\omega_d^{-p/d}} (1 - \omega_d t^{d/p})^m dt = \omega_d^{-p/d} \int_0^1 (1 - u^{d/p})^m du \\ &\geq \omega_d^{-p/d} \int_0^{m^{-p/d}} (1 - u^{d/p})^m du \geq \omega_d^{-p/d} \left(1 - \frac{1}{m}\right)^m m^{-p/d}. \end{aligned}$$

Proof of lower bound

Given any $\sigma \in \mathcal{S}^n$,

$$|X_i - Y_{\sigma(i)}|^p \geq \min_{j=1, \dots, n} |X_i - Y_j|^p,$$

hence

$$B_{\rho, n} \geq \sum_{i=1}^n \min_{j=1, \dots, n} |X_i - Y_j|^p.$$

Taking expectation,

$$\begin{aligned} \mathbb{E}[B_{\rho, n}] &\geq \mathbb{E} \left[\sum_{i=1}^n \min_{j=1, \dots, n} |X_i - Y_j|^p \right] \\ &= \sum_{i=1}^n \mathbb{E} \left[\min_{j=1, \dots, n} |X_i - Y_j|^p \right] \\ &= n \mathbb{E} \left[\min_{j=1, \dots, n} |X_1 - Y_j|^p \right] \quad \text{the law of } (X_i, Y_1, \dots, Y_n) \text{ does not depend on } i \\ &= n \int_{[0,1]^d} \mathbb{E} \left[\min_{j=1, \dots, n} |x - Y_j|^p \right] dx \quad \text{by independence of } X_1 \text{ and } (Y_j)_j \\ &\geq n \cdot c(d, p) n^{-p/d} \quad \text{by Lemma 1.} \end{aligned}$$

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After studying $\mathbb{E}[B_{p,n}]$, it is natural to consider $\text{Var}(B_{p,n})$.

The general phenomenon of **concentration of measure** gives that $B_{p,n}$ is very close to its average because (as described by Talagrand)

the random variable depends in a “smooth” way on a large number of independent variables

Other (better known) examples are

- the law of large numbers, where $\frac{1}{n} \sum_{i=1}^n X_i \rightarrow \mathbb{E}[X_i]$
- Kolmogorov 0-1 laws, yielding e.g. that $\limsup_{n \rightarrow \infty} X_n$ is a.s. constant

Let E be a set. A function $f : E^m \rightarrow \mathbb{R}$ has **bounded differences** if

for every $i \in \{1, \dots, m\}$ there exists $d_i \geq 0$ such that

$$|f(x_1, \dots, x_{i-1}, x_i, x_{i+1}, \dots, x_m) - f(x_1, \dots, x_{i-1}, x'_i, x_{i+1}, \dots, x_m)| \leq d_i$$

for every $x_1, \dots, x_{i-1}, x_i, x_{i+1}, \dots, x_m, x'_i \in E$.

Lemma (Azuma-Hoeffding-McDiarmid)

Let (E, \mathcal{E}) be a measurable space, $f : E^m \rightarrow \mathbb{R}$ with bounded differences and let $(X_i)_{i=1}^m$ be independent random variables with values in E . Then, $Z = f(X_1, \dots, X_m)$ satisfies

$$P(|Z - \mathbb{E}[Z]| > r) \leq 2 \exp\left(-\frac{2r^2}{\sum_{i=1}^m d_i^2}\right) \quad \text{for every } r > 0.$$

Bounded difference as Lipschitz regularity

Consider the distance on E given by $d(x, y) = 1$ if and only if $x \neq y$.

Then, the following are equivalent:

- $f : E^m \rightarrow \mathbb{R}$ has bounded differences
- for every $i = 1, \dots, m$ and $x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_m \in E$, the function

$$x_i \mapsto f(x_1, \dots, x_{i-1}, x_i, x_{i+1}, \dots, x_m)$$

is Lipschitz.

Other concentration inequalities allow for Lipschitz regularity with respect to different distances (e.g., the Euclidean distance on E if $E \subseteq \mathbb{R}^d$).

Lipschitz functions are stable w.r.t. many operations, e.g. **pointwise minima** (with bounded constants).

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Let $E = [0, 1]^d$. For any $\sigma \in \mathcal{S}^n$ we argue that

$$f_\sigma(x_1, \dots, x_n, y_1, \dots, y_n) := \sum_{i=1}^n |x_i - y_{\sigma(i)}|^p$$

has bounded differences (with $d_i = 2d^{p/2}$). The function $\min_{\sigma \in \mathcal{S}^n} f_\sigma$ will also have bounded differences.

Let $i \in \{1, \dots, 2n\}$ and compute the difference of f_σ when evaluated at two points that may differ only for one component:

$$x_i \text{ (if } i \leq n) \text{ or } y_{i-n} \text{ (if } i > n),$$

We find

$$\left| |x_i - y_{\sigma(i)}|^p - |x'_i - y_{\sigma(i)}|^p \right| \leq 2d^{p/2}$$

in the case $i \leq n$ and

$$\left| |x_{\sigma^{-1}(i-n)} - y_{i-n}|^p - |x'_{\sigma^{-1}(i-n)} - y'_{i-n}|^p \right| \leq 2d^{p/2}$$

in the case $i > n$.

We apply the concentration inequality with $m = 2n$ and the variables $(X_1, X_2, \dots, X_n, Y_1, \dots, Y_n)$.

Proposition

For every $r > 0$, one has

$$P(|B_{p,n} - \mathbb{E}[B_{p,n}]| > r) \leq 2 \exp\left(-\frac{r^2}{4d^p n}\right),$$

or equivalently (replacing r with $rn^{1-p/d}$),

$$P\left(\left|\frac{B_{p,n}}{n^{1-p/d}} - \frac{\mathbb{E}[B_{p,n}]}{n^{1-p/d}}\right| > r\right) \leq 2 \exp\left(-\frac{r^2 n^{1-2p/d}}{4d^p}\right).$$

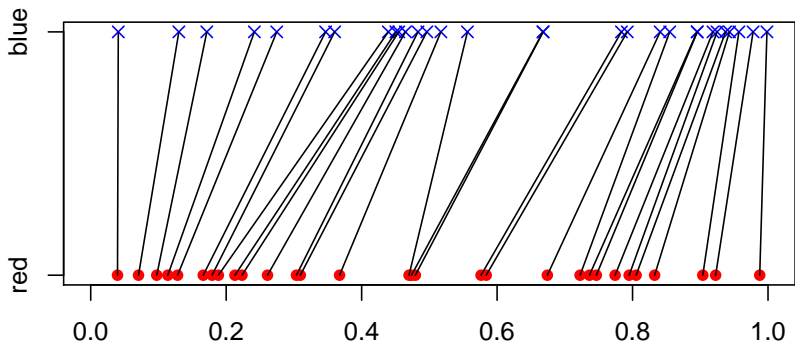
In the last inequality the right hand side is infinitesimal (summable) as $n \rightarrow \infty$ provided that $p < \frac{d}{2}$. Borel-Cantelli lemma yields:

if $p < \frac{d}{2}$ and $\lim_{n \rightarrow \infty} \mathbb{E}[B_{p,n}] n^{1-p/d} = \beta_d$ exists (finite), then $B_{p,n} \rightarrow \beta_d$ P -a.s.

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One-dimensional case

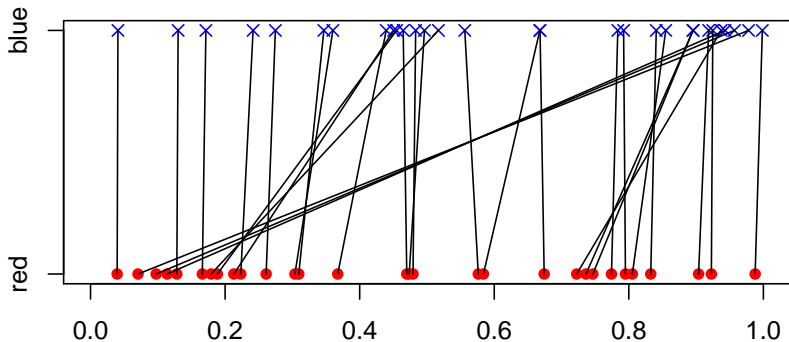
For $d = 1$ and $p \geq 1$, the bipartite matching admits a simple solution.



(Matching of $n = 30$ pairs of i.i.d. uniform points on the unit interval, $p = 1.1$)

One-dimensional case

This is not the case for $p < 1$:



(Matching of $n = 30$ pairs of i.i.d. uniform points on the unit interval, $p = 0.9$)

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Given $n \geq 1$ (distinct) points $(z_i)_{i=1}^n$ in \mathbb{R} , we denote the **k -th smallest** value among them by $z_{(k)}$, e.g.

$$z_{(1)} = \min_i \{z_i\}, \quad z_{(n)} = \max_i \{z_i\},$$

and more generally, for $k \in \{1, \dots, n-1\}$.

$$z_{(k+1)} = \min \{z_i : z_i > z_{(k)}\}.$$

Order statistics

When $z_i = X_i$ are random, write $X_{(k)}$: these are also known order statistics (of the sample).

Proposition

Let $n \geq 1$ and let $(x_i)_{i=1}^n, (y_i)_{i=1}^n$ be distinct points in \mathbb{R} . Then there exists a unique $\sigma^\dagger \in \mathcal{S}^n$ such that

$$\{(x_i, y_{\sigma^\dagger(i)})\}_{i=1}^n = \{(x_{(i)}, y_{(i)})\}_{i=1}^n.$$

Moreover, for every $p \geq 1$,

$$\sum_{i=1}^n |x_i - y_{\sigma^\dagger(i)}|^p = \sum_{i=1}^n |x_{(i)} - y_{(i)}|^p = \min_{\sigma \in \mathcal{S}^n} \sum_{i=1}^n |x_i - y_{\sigma(i)}|^p.$$

Proof **by induction** using the inequality (which is just the case $n = 2$): for any $x < x', y < y'$, then

$$|x - y|^p + |x' - y'|^p \leq |x - y'|^p + |x' - y|^p.$$

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Lemma

Let $(Z_i)_{i=1}^n$ be i.i.d. random variables uniformly distributed on $[0, 1]$. Then for every $t \in [0, 1]$, $k > 1$,

$$P(Z_{(k)} > t > Z_{(k-1)}) = \binom{n}{k-1} t^{k-1} (1-t)^{n-k+1},$$

so that $Z_{(k)}$ has density $\text{Beta}(k, n+1-k)$

$$f_k(t) = \frac{n!}{(k-1)!(n-k)!} t^{k-1} (1-t)^{n-k}.$$

Combining all these facts, we find the following result

Theorem

For $n \geq 1$,

$$\mathbb{E}[B_{2,n}] = \frac{1}{3} \frac{n}{n+1}.$$

In particular,

$$\lim_{n \rightarrow \infty} \mathbb{E}[B_{2,n}] = \frac{1}{3}.$$

This shows that the general lower bound that reads in this case

$$\mathbb{E}[B_{2,n}] \geq c(d, 2)n^{-1},$$

misses in fact the correct order.

As an exercise, try to argue similarly for $B_{4,n}$.

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We describe an approach (due to various authors) to prove existence of

$$\lim_{n \rightarrow \infty} \frac{\mathbb{E}[B_{p,n}]}{n^{1-p/d}} \in (0, \infty),$$

provided that

$$0 < p < \frac{d}{2}$$

Use self-similarity of cube $[0, 1]^d$:

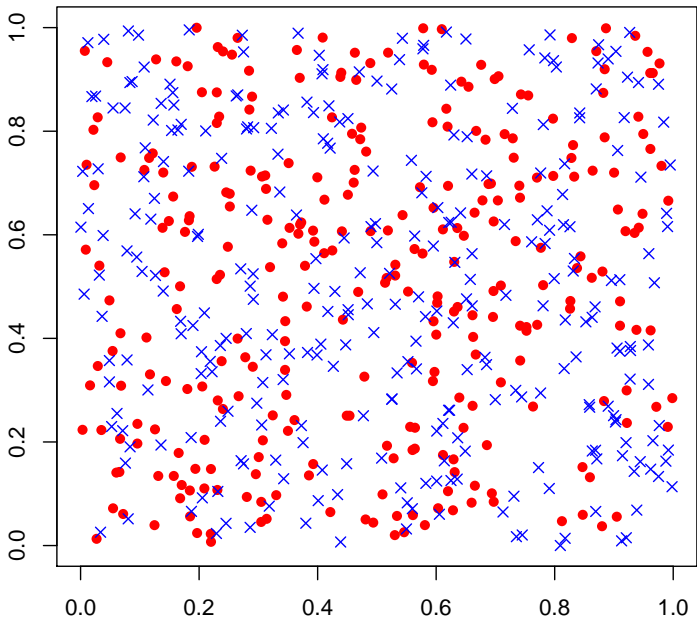
- 1 Decompose it into smaller cubes
- 2 Find a *local matching* on each small cube
- 3 *Glue* them together to find a matching on the initial cube.

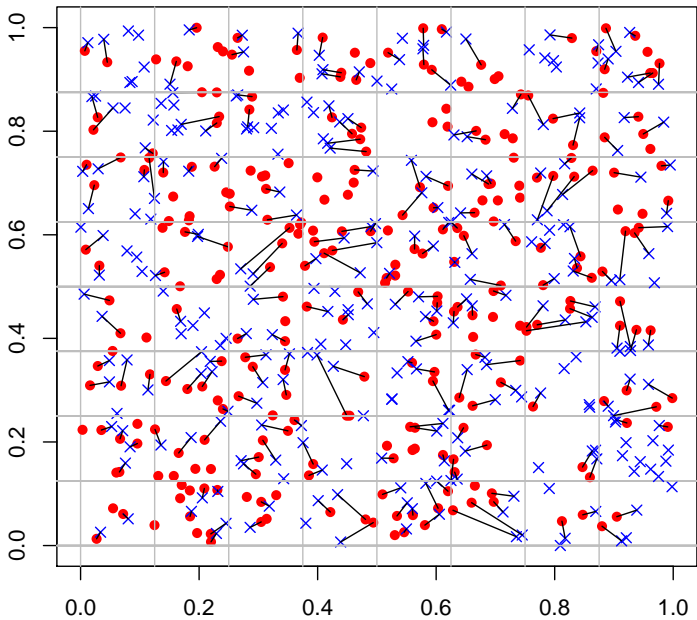
Good: cheaper to compute, one can iterate at multiple scales

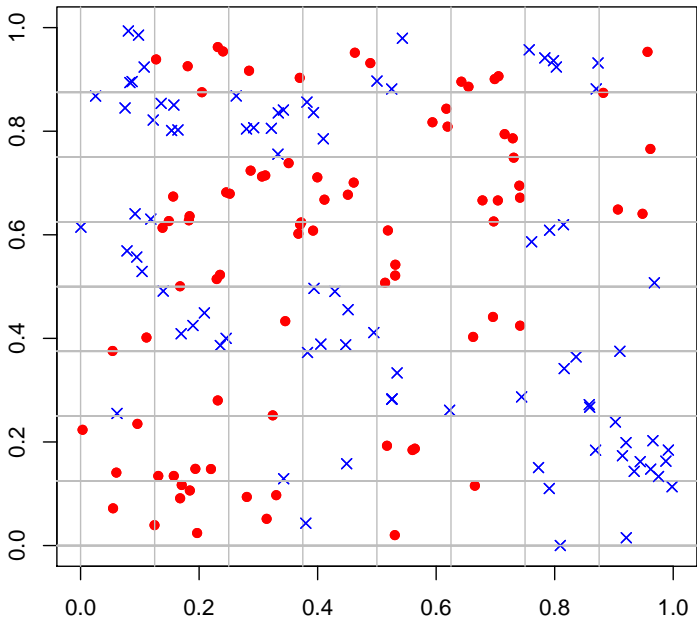
Bad: on smaller cubes the number of red/blue points differ, the matching will not be globally optimal

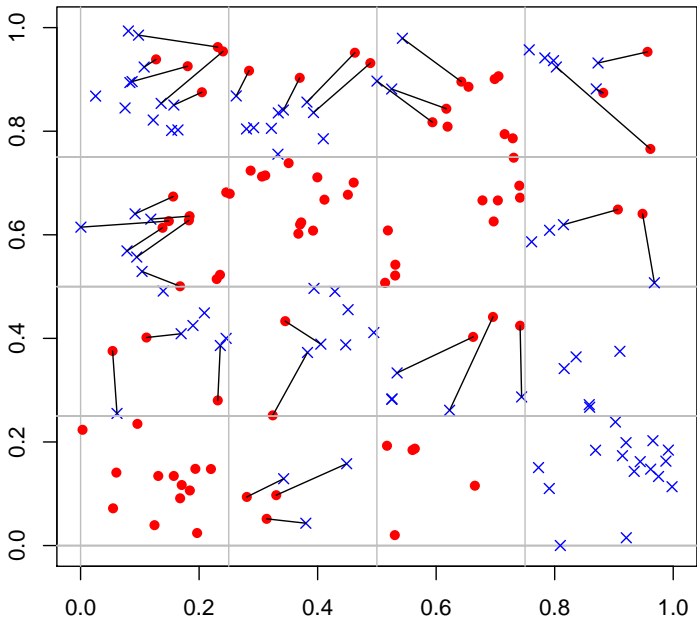
Simulating $n = 300$ pair of points:

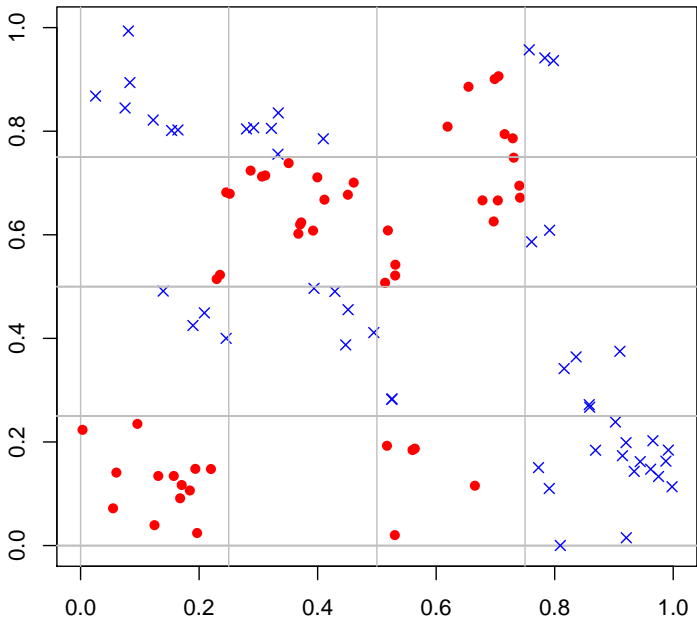
gluing local matchings at scales 2^{-k} with $k = 3, 2, 1, 0$ ($p = 1$)

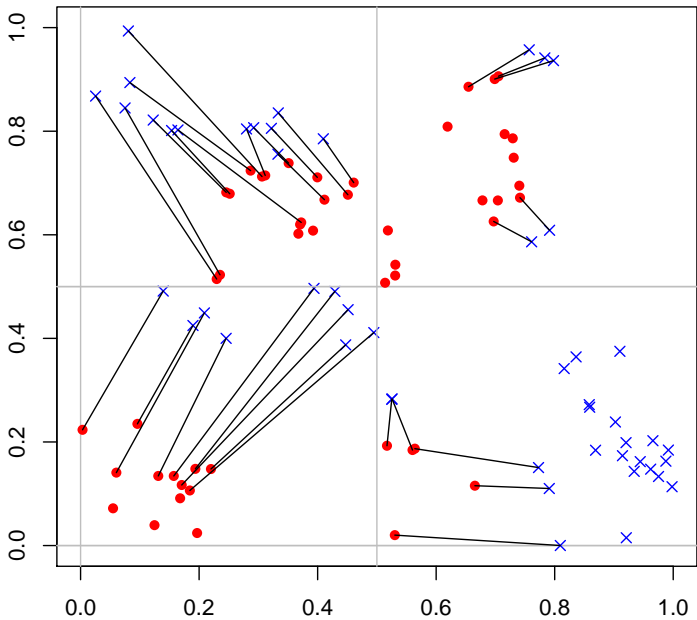


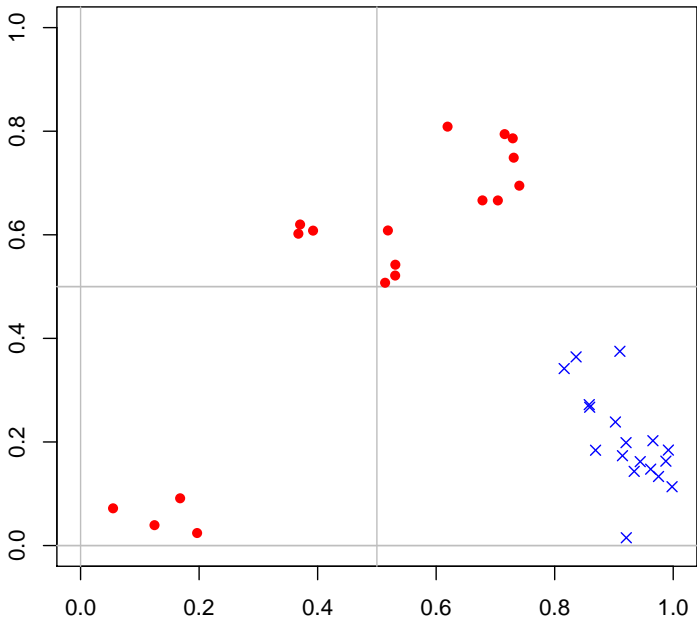


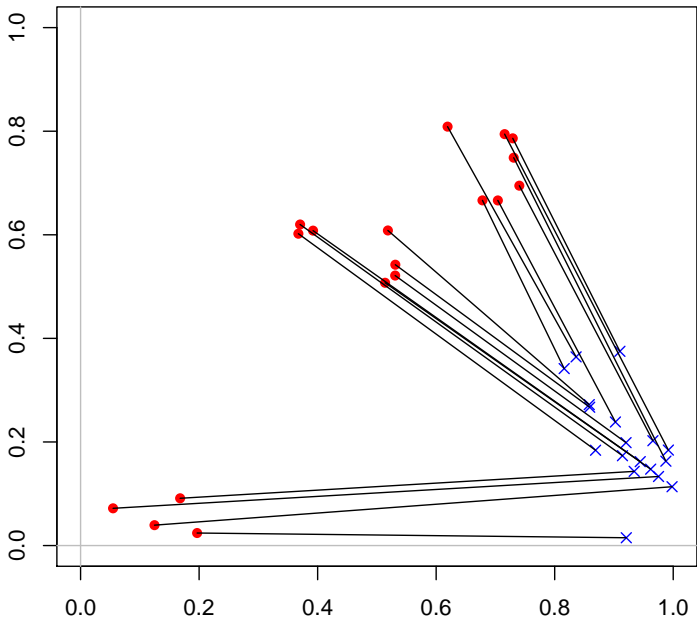












Local matchings: notation

Given $n \in \mathbb{N}$, $x = (x_i)_{i=1}^n \subseteq \mathbb{R}^d$, write

$$\mathcal{I}(x; U) := \{i \in \{1, \dots, n\} : x_i \in U\},$$

$$n(x; U) := \#\mathcal{I}(x; U).$$

For $m \in \mathbb{N}$, $y = (y_i)_{i=1}^m \subseteq \mathbb{R}^d$, a “local” matching in U is a **relation** $\sigma \subseteq \mathcal{I}(x; U) \times \mathcal{I}(y; U)$ **injective functions** from the set with least elements:

- if $n(x; U) \leq n(y; U)$, then $\sigma : \mathcal{I}(x; U) \rightarrow \mathcal{I}(y; U)$ injective
- if $n(x; U) \geq n(y; U)$, then $\sigma : \mathcal{I}(y; U) \rightarrow \mathcal{I}(x; U)$ injective.

Let $\mathcal{S}(x, y; U)$ be set of “local” matchings in U and define the **local cost**

$$b_p(x, y; U) := \min_{\sigma \in \mathcal{S}(x, y; U)} \sum_{(i, j) \in \sigma} |x_i - y_j|^p.$$

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Lemma

Let $n, m \in \mathbb{N}$ $x = (x_i)_{i=1}^n, y = (y_i)_{i=1}^m \subseteq \mathbb{R}^d, U \subseteq \mathbb{R}^d$ Borel. Then,

1 b_p is **local**, i.e., $b_p(x, y; U) = b_p((x_i)_{i \in \mathcal{I}(x; U)}, (y_i)_{i \in \mathcal{I}(y; U)}, U)$.

2 b_p is **translation invariant**: if $v \in \mathbb{R}^d$, then

$$b_p(x + v, y + v; U + v) = b_p(x, y; U),$$

where $x + v = (x_i + v)_{i=1}^n, y + v = (y_i + v)_{i=1}^m, U + v = \{u + v : u \in U\}$.

3 b_p is **p-homogeneous**: if $\lambda \in (0, \infty)$, then

$$b_p(\lambda x, \lambda y, \lambda U) = \lambda^p b_p(x, y; U),$$

where $\lambda x = (\lambda x_i)_{i=1}^n, \lambda y = (\lambda y_i)_{i=1}^m, \lambda U = \{\lambda u : u \in U\}$.

4 b_p is **p-subadditive**: for every (Borel) partition $U = \bigcup_{k=1}^K U_k$, one has

$$b_p(x, y; U) \leq \sum_{k=1}^K b_p(x, y; U_k) + |n(x; U_k) - n(y; U_k)| \text{diam}(U)^p,$$

where $\text{diam}(U) = \sup \{|x - y| : x, y \in U\}$.

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Back to the random case: let $X = (X_i)_{i=1}^n$, $Y = (Y_i)_{i=1}^n$ be i.i.d. uniform on $[0, 1]^d$. Then

$$n(X; U) = \sum_{i=1}^n I_{\{X_i \in U\}}$$

has law $\text{Bin}(n, |U|)$. To obtain a **self-similar** model, we pick n to be also random, with **Poisson law**, with intensity λ , i.e., we impose that

$$P(n(X; [0, 1]^d) = k) = \frac{\lambda^k}{k!} e^{-\lambda} \quad \text{for } k \in \mathbb{N}.$$

Rigorously: fix $\lambda > 0$, let N_X be a Poisson random variables with intensity λ . Let also $(X_i)_{i=1}^{\infty}$ be i.i.d. uniform random variables in $[0, 1]^d$, also independent of N_X . We define a **Poisson point processes of intensity λ** in $[0, 1]^d$ as

$$X := (X_i)_{i=1}^{N_X}$$

(as a random variable with values in $\bigcup_{i=1}^{\infty} \mathbb{R}^{d \times i}$).

Conditionally to $\{N_X = k\}$, then $X = (X_i)_{i=1}^k$ i.i.d. on $[0, 1]^d$ uniformly.

We consider also an independent Poisson point process of same intensity λ

$$Y = (Y_i)_{i=1}^{N_Y}.$$

Lemma

Let $\lambda > 0$ and let $X = (X_i)_{i=1}^{N_X}$ be a Poisson point process of intensity λ on $[0, 1]^d$. For every $U \subseteq [0, 1]^d$ Borel with $|U| > 0$, then

- the random variable

$$n(X; U) = \#\mathcal{I}(X; U)$$

has **Poisson law** of parameter $\lambda|U|$,

- conditionally to $\{n(X; U) = k\}$, the k random variables $(X_i)_{i \in \mathcal{I}(X; U)}$ are independent with uniform law on U .

Corollary

Let $Q = [a, b]^d \subseteq [0, 1]^d$. Then the process

$$((X_i - a)/(b - a))_{i \in \mathcal{I}(X; Q)},$$

is a Poisson point process on $[0, 1]^d$ of intensity $\lambda(b - a)^d$.

When we glue together local matching, we will not find **exact monotonicity** but an approximate one.

Lemma

Let $\alpha > 0$, $c \geq 0$, $f : [1, \infty) \rightarrow [0, \infty)$ be continuous and such that, for every $m \geq 1$ integer, $\lambda \geq 1$ real,

$$f(m\lambda) \leq f(\lambda) + c\lambda^{-\alpha}.$$

Then $\lim_{\lambda \rightarrow \infty} f(\lambda)$ exists.

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Theorem

Let $p \in (0, d/2)$, let X_λ, Y_λ be independent Poisson point processes on $[0, 1]^d$ with intensity λ . Then

$$\lim_{\lambda \rightarrow \infty} \frac{\mathbb{E} [b_p(X_\lambda, Y_\lambda, [0, 1]^d)]}{\lambda^{1-p/d}} \in (0, \infty)$$

exists.

We describe a recent approach to the case $d = 2$: one can show e.g., that

$$\lim_{n \rightarrow \infty} \frac{\mathbb{E}[B_{2,n}]}{\log(n)} = \frac{1}{2\pi}.$$

Again, the rate $\log(n)$ differs from the heuristic $n^{1-p/d} = n^0$. The following problem is open:

$$\text{does } \lim_{n \rightarrow \infty} \frac{\mathbb{E}[B_{1,n}]}{\sqrt{n \log(n)}} \text{ exist?}$$

We focus on the simpler bound, for $p = 1$,

$$\limsup_{n \rightarrow \infty} \frac{\mathbb{E}[B_{1,n}]}{\sqrt{n \log(n)}} < \infty.$$

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The main ideas comes from a more general theory of **optimal transport**, where one looks for efficient **couplings** between probability laws.

Matchings are particular (extreme) cases of couplings.

Analytical tools: **duality formula**

$$B_{1,n} = \max_{f \in \text{Lip}(1)} \left\{ \sum_{i=1}^n f(X_i) - f(Y_i) \right\}$$

where $\text{Lip}(1)$ denotes the set of $f : [0, 1]^d \rightarrow \mathbb{R}$ such that $|f(x) - f(y)| \leq |x - y|$ for every $x, y \in [0, 1]^d$.

Duality for general p

$$B_{p,n} = \max \left\{ \sum_{i=1}^n f(X_i) - g(Y_i) : f(x) - g(y) \leq |x - y|^p \quad \forall x, y \right\}$$

- 1 $(Y_i)_{i=1}^n$ are factories (suppliers)
- 2 $(X_i)_{i=1}^n$ are buyers (cities)
- 3 $|x - y|^p$ is the cost to transport a unit of good from x to y
- 4 A company wants to take care of distribution of goods: buys at price $g(y)$ and sells at price $f(x)$.
- 5 It is competitive if $f(x) - g(y) \leq |x - y|^p$.
- 6 The overall profit is

$$\sum_{i=1}^n f(X_i) - g(Y_i)$$

- 7 Duality reads **max profit = min cost**

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We solve the random elliptic partial differential equation (PDE)

$$-\Delta u = \sum_{i=1}^n (\delta_{X_i} - \delta_{Y_i}) \quad \text{in } [0, 1]^2,$$

where δ_x means the Dirac measure at the point x and use duality (for $p = 1$)

$$\begin{aligned} \sum_{i=1}^n f(X_i) - f(Y_i) &= \int_{[0,1]^2} f \, d \left(\sum_{i=1}^n (\delta_{X_i} - \delta_{Y_i}) \right) \\ &= - \int_{[0,1]^2} f(x) \Delta u(x) \, dx \\ &= \int_{[0,1]^2} \nabla f(x) \nabla u(x) \, dx \leq \int_{[0,1]^2} |\nabla u(x)| \, dx \end{aligned}$$

Taking expectation, we obtain the inequality

$$\mathbb{E} [B_{1,n}] \leq \mathbb{E} \left[\int_{[0,1]^2} |\nabla u(x)| \, dx \right] \leq \mathbb{E} \left[\int_{[0,1]^2} |\nabla u(x)|^2 \, dx \right]^{1/2}.$$

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The advantage of the PDE with respect to the matching problem is **linearity**: explicit formulas e.g. via Fourier series:

$$u(x_1, x_2) = 4 \sum_{m \in \mathbb{N}^2} a_m \cos(m_1 \pi x_1) \cos(m_2 \pi x_2),$$

where $m = (m_1, m_2)$, $|m| = (m_1^2 + m_2^2)^{1/2}$, $a_{(0,0)} = 0$ and

$$a_m = \frac{4}{\pi^2 |m|^2} \sum_{i=1}^n \cos(m_1 \pi X_{i,1}) \cos(m_2 \pi X_{i,2}) - \cos(m_1 \pi Y_{i,1}) \cos(m_2 \pi Y_{i,2}).$$

Fourier coefficients are random variables, with mean 0 and variance

$$\mathbb{E} [a_m^2] = \frac{2n}{(\pi^2 |m|^2)^2}.$$

Recall

$$\mathbb{E} \left[\mathbf{a}_m^2 \right] = \frac{2n}{\left(\pi^2 |m|^2 \right)^2}.$$

Plancherel identity gives

$$\int_{[0,1]^2} |\nabla u(x)|^2 dx = \sum_{m \in \mathbb{N}^2} \pi^2 |m|^2 \mathbf{a}_m^2.$$

Taking expectation we obtain the series

$$\begin{aligned} \mathbb{E} \left[\int_{[0,1]^2} |\nabla u(x)|^2 dx \right] &= \sum_{m \in \mathbb{N}^2} \pi^2 |m|^2 \mathbb{E} \left[\mathbf{a}_m^2 \right] \\ &= \frac{2n}{\pi^2} \sum_{\substack{m \in \mathbb{N}^2 \\ m \neq (0,0)}} \frac{1}{|m|^2} \\ &\approx \frac{2n}{\pi^2} \int_1^\infty \frac{1}{r} dr \frac{\pi}{2} = \infty! \end{aligned}$$

However, partial sums

$$\sum_{1 \leq |m| \leq a} \frac{1}{|m|^2} \approx \frac{\pi}{2} \int_1^a r \, dr = \frac{\pi}{2} \log(a)$$

CLAIM: Only frequencies m with $|m| \leq n$ matter \Rightarrow

$$\mathbb{E} [B_{1,n}] \leq c(d) \sqrt{n \log(n)}.$$

Difficult to make rigorous!