# Un invito al Trasporto Ottimo Quantistico<sup>1</sup>

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XXII Congresso UMI September 8, 2023

<sup>1</sup>Based on joint works with G. De Palma, M. Marvian, S. Lloyd, T. Titkos and D. Virosztek

QOT arXiv:2307.16268

## Outline

- Classical Optimal Transport
- Quantum Systems
- Quantum Optimal Transport



## Plan

Classical Optimal Transport

- Monge
- Kantorovich
- Earth Mover's distance

#### Quantum Systems

- 3 Quantum Optimal Transport
- 4 Conclusion

# Monge's transport problem

#### Monge (1781): sur la théorie des déblais et des remblais.



#### How to transport soil during a construction with minimal expenses?

# The assignment problem

A discrete formulation: given a

• cost c(x, y) of moving unit of soil from position x to position y, e.g.

 $\boldsymbol{c}(\boldsymbol{x},\boldsymbol{y})=|\boldsymbol{x}-\boldsymbol{y}|,$ 

- Source distribution of soil  $\sigma = (\sigma(x_i))_i$
- Target distribution (dump)  $\rho = (\rho(y_j))_j$

Find  $T : \{x_i\} \to \{y_j\}$  that moves  $\sigma$  into  $\rho$  with minimal transport cost

$$\sum_i c(x_i, T(x_i))\sigma(x_i).$$

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# Kantovorich and linear programming

Relax the map T to a transport plan

$$T(x_i, y_j) \geq 0$$

such that

$$\sum_{j} T(x_i, y_j) = \sigma(x_i), \quad \sum_{i} T(x_i, y_j) = \rho(y_j).$$

Probabilistic intepretation:

$$K(y_j|x_i) := \frac{T(x_i, y_j)}{\sigma(x_i)} \in [0, 1].$$

The variational problem becomes

$$\min_{T} \sum_{i} \sum_{j} c(x_i, y_j) T(x_i, y_j)$$

 $\Rightarrow$  linear programming!

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# From soil to probabilities

#### • For a transport plan T it must be

$$\sum_{i} \sigma(\mathbf{x}_{i}) = \sum_{i} \sum_{j} T(\mathbf{x}_{i}, \mathbf{y}_{j}) = \sum_{j} \sum_{i} T(\mathbf{x}_{i}, \mathbf{y}_{j}) = \sum_{j} \rho(\mathbf{y}_{j}).$$

• We assume that  $\rho$ ,  $\sigma$  are probability mass functions (discrete densities):

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## Earth Mover's distance and duality

• If c(x, y) = d(x, y) is a distance, then

$$W_1(\sigma,\rho) = \min_T \sum_i \sum_j d(x_i, y_j) T(x_i, y_j)$$

#### induces a distance (Wasserstein-Kantorovich, Earth Mover's).

• Kantorovich duality:

$$W_1(\sigma,\rho) = \max\left\{\sum_i f(x_i)\sigma(x_i) - \sum_j f(y_j)\rho(y_j) : |f(x) - f(y)| \le d(x,y)\right\}.$$

• Buy at price  $f(y_i)$  and sell at price  $f(x_i) \Rightarrow$  maximize the profit!

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# $c(x, y) = |x - y|^{p}, p = 0.6$



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- geometric interpolation between distributions
- Idiscriminator in generative AI models (WGANs)
- functional inequalities (isoperimetric, concentration of measure)
- PDE's as gradient flows
- geometry (synthetic Ricci curvature bounds)

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## Plan



2

Quantum SystemsFrom Classical to Quantum

- Systems of qubits
- 3 Quantum Optimal Transport

#### 4 Conclusion

#### Classical



#### Quantum



 $\sum_{x\in E} f(x)$ 

#### H Hilbert space (C') : H Hilbert space

A : A --> A linear self-adjoint (observable) non-negative 14/<sup>2</sup> -- A A

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$f: E \to \mathbb{C}$ real-valued non-negative $ f ^2$	
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angle\in H$ V < H

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Tr[*A*]

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Markov (transition) operator  $K : \mathcal{P}(E) \rightarrow \mathcal{P}(E)$  $(Kp)(y) = \sum_{x \in E} K(y|x)p(x)$ 

Shannon's entropy  $S(p) = -\sum_{x} p(x) \log p(x)$ 

Relative entropy (KL divergence)  $D(p||q) = \sum_{x} p(x) \log(p(x)/q(x))$   $\varphi \in \mathcal{S}(H)$  , states  $\varphi \in \mathcal{S}(H)$  , states  $\varphi = |\psi\rangle\langle\psi|$ 

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Von Neumann entropy  $S(\rho) = - \operatorname{Tr}[\rho \log \rho]$ 

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# Single qubit system

A quantum analogue of  $\{0,1\}.$  Set

$$H = \mathbb{C}^2$$
.

Standard (computational) basis

$$\{|0\rangle, |1\rangle\} = \{(1,0), (0,1)\},\$$

Pauli operators

$$\sigma_{x} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad \sigma_{y} = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, \quad \sigma_{z} = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}.$$

Parametrization of states:

$$\rho = \frac{1}{2} \left( \mathbb{I}_{\mathbb{C}^2} + b_x \sigma_x + b_y \sigma_y + b_z \sigma_z \right),$$

with

$$(b_x, b_y, b_z) \in \mathbb{R}^3, \quad b_x^2 + b_y^2 + b_z^2 \leq 1.$$

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$$\rho = |\psi\rangle\langle\psi|$$
 is pure  $\Leftrightarrow$   $b_x^2 + b_y^2 + b_z^2 = 1$ .



Figure: The Bloch representation of a pure state  $\rho = |\psi\rangle\langle\psi| \in \mathbb{C}^2$ .

# Many-qubits systems

A quantum analogue of  $\{0, 1\}^n$ :

$$H_n = (\mathbb{C}^2)^{\otimes n} \sim \mathbb{C}^{2^n}.$$

Standard (computational) basis

$$\{|\boldsymbol{s}\rangle : \boldsymbol{s} \in \{0,1\}^n\}.$$

Pure states are  $\rho = |\psi\rangle\langle\psi|$  but not necessarily  $\psi \in \{0, 1\}^n$ :

$$ert \psi 
angle = rac{1}{2^{n/2}} \sum_{s \in \{0,1\}^n} ert s 
angle$$
 (uniform superposition  
 $ert \Phi^+ 
angle = rac{ert 00 
angle + ert 11 
angle}{\sqrt{2}}$  (Bell state,  $n = 2$ )

"Classical" probabilities on  $\{0,1\}^n$  are diagonal in the computational basis:

$$o = \sum_{\boldsymbol{s} \in \{0,1\}^n} \boldsymbol{p}(\boldsymbol{s}) | \boldsymbol{s} \rangle \langle \boldsymbol{s} |.$$

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### Plan

Classical Optimal Transport

- Quantum Systems
- Quantum Optimal Transport
  - Overview
  - Transport via quantum channels
  - Quantum earth mover's distance

### Conclusion



### Classical distances between probabilities have quantum analogues:

- Total variation  $\rightarrow$  Trace distance
- Hellinger distance → Fidelity
- KL divergence → Relative entropy

Like their classical counterparts:

- Quite general, easy to compute or approximate
- Not adapted to specific geometry.

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- 2013 Agredo: 1-Wasserstein extending any distance on basis vectors
- 2016 Golse/Mouhot/Paul: guantum Kantorovich problem

 De Palma/T:: Guantum optimal transport using channels

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- 2019 De Palma/T.: Quantum optimal transport using channels
- 2020 De Palma/Marvian/T./Lloyd: Earth mover's distance on qubits

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#### Overview

### A timeline

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• Any Kantorovich transport plan  $T(x_i, y_j)$  yields a transition kernel

$$\mathcal{K}(\mathbf{y}_{i}|\mathbf{x}_{i}) = \frac{\mathcal{T}(\mathbf{x}_{i},\mathbf{y}_{j})}{\sigma(\mathbf{x}_{i})},$$

such that

$$K\sigma = \rho$$

• Fix real-valued functions  $(g_\ell)_{\ell=1}^d$  and choose a "quadratic" cost

$$c(x,y) = \sum_{\ell=1}^{d} (g_{\ell}(x) - g_{\ell}(y))^2.$$

$$\sum_{i} \sum_{j} c(x_i, y_j) T(x_i, y_j)$$
  
=  $\sum_{\ell} \sum_{i} g_{\ell}^2(x_i) \sigma(x_i) + \sum_{i} g_{\ell}^2(y_j) \rho(y_j) - 2 \sum_{i} g_{\ell}(x_i) (K^{\dagger} g_{\ell})(x_i) \sigma(x_i),$ 

where

$$(K^{\dagger}g)(x_i) = \sum g(y_j)K(y_j|x_i)$$

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• Fix real-valued functions  $(g_\ell)_{\ell=1}^d$  and choose a "quadratic" cost

$$c(x,y)=\sum_{\ell=1}^d(g_\ell(x)-g_\ell(y))^2.$$

$$\sum_{i} \sum_{j} c(x_i, y_j) T(x_i, y_j)$$
  
=  $\sum_{\ell} \sum_{i} g_{\ell}^2(x_i) \sigma(x_i) + \sum_{j} g_{\ell}^2(y_j) \rho(y_j) - 2 \sum_{i} g_{\ell}(x_i) (\mathcal{K}^{\dagger} g_{\ell})(x_i) \sigma(x_i),$ 

where

$$(K^{\dagger}g)(x_i) = \sum g(y_j)K(y_j|x_i)$$

Dario Trevisan (UNIPI)

• Given  $\rho$ ,  $\sigma \in S(H)$ , define transport plans as quantum channels

$$\Phi: \mathcal{S}(\mathcal{H}) \to \mathcal{S}(\mathcal{H}), \quad \Phi(\tau) = \sum_{j} B_{j} \tau B_{j}^{\dagger}$$

such that

$$\Phi(\sigma) = \rho.$$

• Fix observables  $\mathcal{R} = (R_{\ell})_{\ell=1}^{d}$  and define  $C_{\mathcal{R}}(\Phi) = \sum_{\ell} \operatorname{Tr}[R_{\ell}^2 \sigma] + \operatorname{Tr}[R_{\ell}^2 \rho] - 2 \operatorname{Tr}[R_{\ell} \sqrt{\sigma}(\Phi^{\dagger} R_{\ell}) \sqrt{\sigma}],$ 

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$$\Phi^{\dagger}R = \sum_{j} B_{j}^{\dagger}RB_{j}$$

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 Aim: define a quantum analogue of earth mover's distance with respect to the Hamming distance on binary strings:

$$|x - y|_{\mathsf{H}} = \sum_{i=1}^{n} |x_i - y_i|$$
 for  $x, y \in \{0, 1\}^n$ .

• Recall the Kantorovich duality:

$$W_{1}(\sigma,\rho) = \max\left\{\sum_{x \in \{0,1\}^{n}} f(x)(\sigma(x) - \rho(x)) : |f(x) - f(y)| \le |x - y|_{H}\right\}$$

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• Equivalently:

 $\|f\|_{Lip} = 2 \max_{i=1,\dots,n} \min \left\{ \|f - g_i\|_{\infty} : g_i : \{0,1\}^n \to \mathbb{R} \text{ independent of site } i \right\}$ 

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## Plan

Classical Optimal Transport

Quantum Systems

3 Quantum Optimal Transport





Quantum OT has already found applications:

- Non-commutative Ricci curvature bounds
- Quantum machine learning (generative models, quantum neural nets)
- Equilibrium measures on quantum spin systems

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- Geometry induced by OT (isometries, inequalities)
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- Structural properties of optimizers
- Connections among different approaches...

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