

Logarithmic geometry and tropicalization ← "piecewise linear" AG

"decorated" version of algebraic geometry,
where varieties are equipped with

a sheaf of monoids \rightsquigarrow combinatorics,
polyhedral geometry
↑
commutative

Plan: 1. Toric varieties

2. Log geometry

3. Tropical geometry and tropicalization.

1. Toric varieties [Cox-Little-Schenck, Toric varieties]

Def: a monoid is a set P with $+$: $P \times P \rightarrow P$ associative and commutative, with a neutral elt. $0 \in P$.

Ex: • every abelian gp is a monoid

• \mathbb{N} with $+$ and 0 is a monoid, \mathbb{N}^k with component-wise $+$ and $(0, \dots, 0)$.

(products $P \times Q$ are as you'd expect)

• $P = \{0, 1\}$, with $1 + 1 = 1$

Def: a homomorphism of monoids $f: P \rightarrow Q$ is

a function compatible with + and 0.

Associated group: $P \rightsquigarrow P^{\text{gp}} = \frac{P \times P}{\sim}$

idea: $(p, q) \leftrightarrow p - q$

$$(p, q) \sim (p', q') \Leftrightarrow p + q' + r = p' + q + r \\ (\exists r \in P \text{ s.t.})$$

This becomes a group with the induced operation.

There is a homom. $P \rightarrow P^{\text{gp}}$ and this is universal,
 $p \mapsto [(p, 0)]$

in the sense that if other ab. group G , every hom.

$P \rightarrow G$ will factor.
 $\downarrow P^{\text{gp}} \exists!$

$(-)^{\text{gp}}: (\text{Mon}) \rightarrow (\text{Grp})$ is the left adjoint of the inclusion

$$(\text{Grp}) \hookrightarrow (\text{Mon}) \quad \text{Hom}_{\text{Ab}}(P^{\text{gp}}, G) = \text{Hom}_{\text{Mon}}(P, G)$$

Def: a monoid is integral if $p+r=q+r$ in $P \Rightarrow p=q$

(equivalent to saying that $P \rightarrow P^{\text{gp}}$ is injective)

and saturated if it is integral, and

$$x \in P^{\text{gp}}, n x \in P \Rightarrow x \in P$$

$n \in \mathbb{N} \quad n \neq 0$

Ex: • \mathbb{N}^k is saturated, $(\mathbb{N}^k)^{\text{gp}} = \mathbb{Z}^k$

• $P = \langle 2, 3 \rangle \subseteq \mathbb{N}$

$P^{\text{gp}} = \mathbb{Z}$

" $\{a \cdot 2 + b \cdot 3 \mid a, b \in \mathbb{N}\}$

$1 = 3 - 2$

but $1 \notin P$, so P is not saturated.

• $P = \{0, 1\}$ with $1+1=1$ is not even integral.

Def: a monoid is finitely generated if there is a surjection $\mathbb{N}^k \twoheadrightarrow P$ for some $k \in \mathbb{N}$

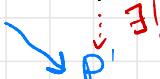
$$(\exists p_1, \dots, p_k \in P \text{ s.t. } p \in P \rightsquigarrow p = \sum_{i=1}^k a_i p_i)$$

fine = integral and finitely generated.

fs = fine and saturated

A presentation of a monoid is a coequalizer diagram

$$\mathbb{N}^k \begin{array}{c} \xrightarrow{g} \\ \xrightarrow{f} \\ \xrightarrow{h} \end{array} \mathbb{N}^n \xrightarrow{f} P$$



$$P = \langle \text{generators} \mid \text{relations} \rangle$$

$f(e_i) \quad g(e_j) = h(e_j)$

Thm (Redei): every finitely generated monoid is finitely presented.

Ex.: $P = \langle p, q, r \mid p+q = 2r \rangle \cong \langle (1,0), (1,2), (1,1) \rangle \subseteq \mathbb{Z}^2$

$P = \langle p, q, r, s \mid p+q = r+s \rangle \cong \langle (1,0,0), (0,1,0), (0,0,1), (1,1,-1) \rangle$

Def.: a monoid is toric if it is fs + torsion free

$k = \mathbb{K}$, think
 $k = \mathbb{C}$ if you want

i.e. $nx = ny \Rightarrow x = y$
 $n \in \mathbb{N}, x, y \in P$
 $n \neq 0$

Def.: the monoid algebra of P is $k[P] = \bigoplus_{p \in P} k \cdot t^p$

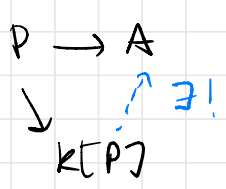
is the free k -vect. space
 on the set P

with the multiplication defined by $t^p \cdot t^q = t^{p+q}$,
 extended by linearity.

Every k -algebra gives a monoid using its multiplication,

and the obvious $P \rightarrow k[P]$ (hom. of monoids)
 $p \mapsto t^p$

is initial w.r.t. hom $P \rightarrow A$ ← any k -algebra



(i.e. the functor $P \mapsto k[P]$ is
 the left adjoint to the forgetful)

In particular a presentation $\mathbb{N}^k \xrightarrow{g} \mathbb{N}^n \rightarrow P$

will give a coequalizer in k -algebras

$$k[\mathbb{N}^k] \xrightarrow{g} k[\mathbb{N}^n] \rightarrow k[P]$$

" "
" "
" "

$$k[y_1, \dots, y_k] \quad k[x_1, \dots, x_n]$$

$$\Rightarrow k[P] \cong \frac{k[x_1, \dots, x_n]}{(g(y_i) - h(y_i))}$$

Ex: $P = \langle p, q, r \mid p+q = 2r \rangle \rightsquigarrow k[P] = \frac{k[x, y, z]}{xy - z^2}$

$$P = \langle p, q, r, s \mid p+q = r+s \rangle \rightsquigarrow k[P] = \frac{k[x, y, z, w]}{(xy - zw)}$$

From now on P will be toric, in part.

$$k[P] \cong \frac{k[x_1, \dots, x_n]}{I}$$

$k[P]$ has an associated affine dg-variety,

$$\text{Spec}(k[P]) = V(I) \subseteq \mathbb{A}_k^n \leftarrow \begin{array}{l} n \text{ dim'l affine space} \\ \uparrow \\ \text{vanishing locus} \end{array}$$

($X = \text{Spec} A \Leftrightarrow X$ is affine, and $A = k[X]$ is its coordinate ring)
 (normal)

Def: an affine toric variety (TV) is an affine var. / k with a dense open immersion $T \subseteq X$, where

T is an alg. forms, and the translation action

$$T \times T \rightarrow T \quad \text{extends to an action } T \times X \rightarrow X.$$
$$(x, y) \mapsto x \cdot y$$

Thm: $A_p := \text{Spec } k[P]$ is an affine TV, and all affine TVs arise like this.

Where is the locus of A_p ? $p \in P^{\text{gp}} \rightsquigarrow k[P] \subseteq k[P^{\text{gp}}]$

$P^{\text{gp}} \cong \mathbb{Z}^r$ for some $r \in \mathbb{N}$, and

$$k[P^{\text{gp}}] \cong k[\mathbb{Z}^r] = k[x_1^{\pm 1}, \dots, x_r^{\pm 1}]$$

$$\text{Spec } k[x_1^{\pm 1}, \dots, x_r^{\pm 1}] = \left(\text{Spec } k[t, t^{-1}] \right)^r \cong (k^*)^r$$

$$k[x_1^{\pm 1}, \dots, x_r^{\pm 1}] = k[x_1^{\pm 1}] \otimes_k k[x_2^{\pm 1}] \otimes_k \dots \otimes_k k[x_r^{\pm 1}]$$

$$\text{Spec } k[t, t^{-1}] \cong \mathbb{A}^1 \setminus \{0\} \quad (k^*)$$

A (split) algebraic torus is a group variety isomorphic to $(k^*)^r$.

Since you can realize $k[P] \subseteq k[P^{\text{gp}}]$ by localizing at

$\prod t^{p_i}$ where $\{p_i\}$ is a generating set of P ,

$T_p = \text{Spec } k[P^{\text{gp}}] \subseteq \text{Spec } k[P]$ is a principal open.

The action $T_P \times \mathbb{A}^p \rightarrow \mathbb{A}^p$ is given by

$$K[P] \rightarrow K[P \otimes P] \otimes_{K[P]} K[P]$$

$$t^P \longmapsto t^P \otimes t^P$$

Ex: $P = \langle (1,0), (1,2), (1,1) \rangle \subseteq \mathbb{Z}^2$, $P^\otimes = \mathbb{Z}^2$

$$K[P] = K[x,y,z] / \langle xy - z^2 \rangle \subseteq K[P^\otimes] = K[s, t^{\pm 1}]$$

$$s \leftrightarrow (1,0)$$

$$t \leftrightarrow (0,1)$$

$$(1,0) \leftarrow x \longmapsto s$$

$$(1,2) \leftarrow y \longmapsto st^2$$

$$(1,1) \leftarrow z \longmapsto st$$

so that $T_P \times \mathbb{A}^p \rightarrow \mathbb{A}^p$ looks like

$$(s,t) \cdot (a,b,c) = (sa, st^2b, stc)$$

$$ab = c^2$$

Remark: normality of $\mathbb{A}^p \iff$ saturatedness of P

Ex: $\text{Spec } K[\langle 2,3 \rangle] \cong \text{Spec } K[x,y] / \langle x^3 - y^2 \rangle$, a non-normal

affine TV

(normal)



and integral
red, sep. scheme of finite
type over K

Def: a toric variety is a normal variety, w/ a dense open
torus $T \subseteq X$ s.t. the translation action of

The dual of $\sigma \subseteq N_{\mathbb{R}}$ is $\sigma^{\vee} = \{ m \in M_{\mathbb{R}} \mid \langle m, n \rangle \geq 0 \ \forall n \in \sigma \}$
 $\subseteq M_{\mathbb{R}}$

This is again a RPC, and $(\sigma^{\vee})^{\vee} = \sigma$.

The point is that every toric monoid P can be written

as $\sigma^{\vee} \cap M$ for a strongly convex RPC in $N_{\mathbb{R}}$

*→ 0-dim'l
doesn't contain any lin. subspace.*

Ex: $P = \langle (1,0), (1,2), (1,1) \rangle = \sigma^{\vee} \cap M$

where $\sigma = \text{Cone}((0,1), (2,-1)) \subseteq \mathbb{R}^2$



Def: a fan is a finite collection of strongly convex RPC

$\Sigma = \{\sigma\}$ s.t.

- $\sigma \cap \sigma'$ is a face of both σ and σ' $\forall \sigma, \sigma' \in \Sigma$.

(a face of σ , $\tau \leq \sigma$, is the intersection of σ with a "supporting hyperplane")

- if $\tau \leq \sigma$, $\sigma \in \Sigma \Rightarrow \tau \in \Sigma$

From a fan Σ in $N_{\mathbb{R}}$, get a TV X_{Σ} :

- $\sigma \in \Sigma \rightsquigarrow X_{\sigma} := \mathbb{A}^p = \text{Spec } k[P]$, $P = \sigma^{\vee} \cap M$
- $\tau \leq \sigma \rightsquigarrow X_{\tau} \hookrightarrow X_{\sigma}$ an open embedding.
- $\sigma \cap \sigma' = \tau \rightsquigarrow X_{\sigma}, X_{\sigma'}$ can be glued along X_{τ} , a common open subvariety.
- do all the gluings prescribed by the fan, to get X_{Σ} .

Thm: X_{Σ} is a TV and every TV arises like this.

Ex: • toric curves:





k^*


\mathbb{A}^1

$$\mathbb{A}^1 \cup \mathbb{A}^1 = \mathbb{P}^1$$

$$\sigma_1(k^*)\sigma_2$$

-  $\mathbb{R}_{\geq 0}^2 = \sigma \rightsquigarrow \mathbb{A}^2$

-  $\rightsquigarrow \mathbb{P}^2$ (\mathbb{P}^n has a similar higher dim'l fan)

-  $\rightsquigarrow \mathbb{P}^1 \times \mathbb{P}^1$ ($X_{\Sigma} \times X_{\Sigma'} \cong X_{\Sigma \times \Sigma'}$)
in general

Slogan: the (equivariant) geometry of X_{Σ} is controlled

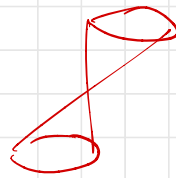
by the "combinatorics" of Σ .

- E_x :
- X_Σ is proper (\Leftrightarrow compact) $\Leftrightarrow |\Sigma| = \bigcup_{\sigma \in \Sigma} \sigma = N_{\mathbb{R}}$
 - X_Σ is smooth \Leftrightarrow every $\sigma \in \Sigma$ is smooth, i.e. the minimal ray generators are part of a \mathbb{Z} -basis of N .

for example, the cone $\sigma = \text{cone}((0,1), (2,-1))$ of yesterday is not smooth, and the corresponding

$X_\sigma = \text{Spec } k[x,y,z] / (xy - z^2)$ has singularity at $(0,0,0)$

"quadric cone"



(real trace)

You can show that σ is smooth $\Leftrightarrow X_\sigma \cong \mathbb{A}^k \times (\mathbb{A}^*)^{n-k}$.

This implies that smooth toric varieties are covered

by affine spaces, and that $X_\Sigma \setminus T$, the "toric boundary" is a simple normal crossings divisor

↑
irred. components
are smooth

↑
the components
intersect transversely

- π_1 , divisors/line bundles, (co)homology, Chow ring, equivariant vector bundles, ... of X_Σ .

can all be described / computed from the fan Σ

Remark: projective TVs are encoded by lattice polytopes, via the "normal fan" construction.

Orbit-cone correspondence

there is a 1:1 correspondence
inclusion-reversing

$\Sigma \leftrightarrow$ orbit closures for the action of T on X_Σ

(\leftrightarrow orbits of the action)

$\sigma \leftrightarrow V(\sigma) = \overline{O(\sigma)} \subseteq X_\Sigma$
($O(\sigma)$)

\simeq torus

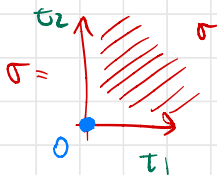
$V(\sigma)$ is actually a $O(\sigma)$ -toric variety.

Ex: $X_\Sigma = \mathbb{A}^2 \cong (\mathbb{K}^*)^2$

$(a,b) \cdot (x,y) = (ax, by)$

\uparrow
 $(\mathbb{K}^*)^2$

orbits: $\{(0,0)\}$, x -axis (origin, τ_2), y -axis (origin, τ_1), $(\mathbb{K}^*)^2$



Remark: the boundary $X_\Sigma \setminus T$ is a union of toric varieties of lower dimension, the $V(\sigma)$ for $\sigma \neq 0$

logarithmic geometry

[Abramovich et al - logarithmic geometry and moduli]

The geometry of a TV is quite constrained, for example
TVs are all rational (birational to \mathbb{P}^n).

log structures aim to generalize the theory of TVs by
retaining some of the combinatorial aspects, but on more
complicated varieties.

$\approx \text{Spec } k[t^{\pm 1}]$

Recall that on an affine TV \mathbb{A}^1/k we have $P \rightarrow k[P]$

$$p \mapsto t^p$$

$$p+t^p \mapsto t^p \cdot t^{p^1}$$

"exponential"

Idea: globalize this homomorphism by "sheafifying".

Def: a log structure on a scheme X

(think algebraic variety, or even
complex analytic space)

is (M, α) where:

pre-log
structure

- M is a sheaf of monoids (for the étale topology,
think Euclidean top. $/\mathbb{C}$)

- $\alpha: M \rightarrow (\mathcal{O}_X, \cdot)$

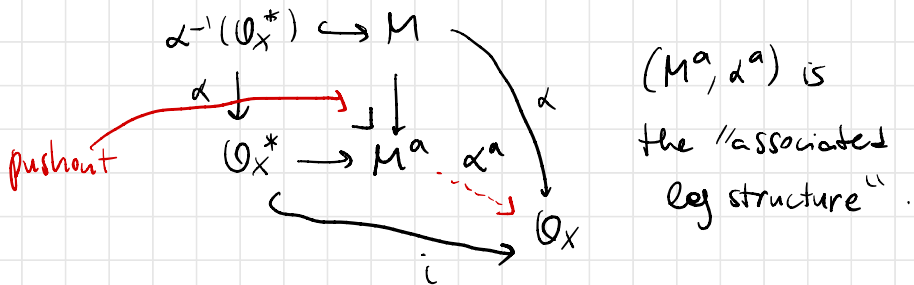
↖ sheaf of regular functions, i.e. morphisms to \mathbb{A}^1

is a homom. of monoids

- $\alpha|_{\alpha^{-1}(\mathcal{O}_X^*)} : \alpha^{-1}(\mathcal{O}_X^*) \xrightarrow{\cong} \mathcal{O}_X^*$ is an isomorphism.

A log scheme is a (X, M, α) of a scheme w/ a log str.

Remark: every pre-log str. can be canonically turned into a log str, as follows: given $\alpha: M \rightarrow \mathcal{O}_X$



A morphism $(M, \alpha) \rightarrow (M', \alpha')$ is a homom. $M \rightarrow M'$ s.t.

$M \rightarrow M'$ commutes.

$$\begin{array}{ccc} \alpha \downarrow & & \downarrow \alpha' \\ & & \mathcal{O}_X \end{array}$$

Ex: • $M = \mathcal{O}_X^*$ and $\alpha = i: \mathcal{O}_X^* \hookrightarrow \mathcal{O}_X$ is the trivial

log structure on X .

simple normal crossing

• if (X, D) is a SNC pair, then one can construct

the divisorial a log str. as follows

log str. ind. by D

$$M_D(U) = \left\{ f \in \mathcal{O}_X(U) \mid f|_{(X, D) \cap U} \in \mathcal{O}^* \right\} \xrightarrow{\alpha} \mathcal{O}_X(U)$$

$U \subseteq X$

~ local equations for the components of D.



If at $x \in X$, $f_1 = 0, \dots, f_k = 0$ are local eqns around x of the components of D , then

$$M_{D, X} \cong \mathcal{O}_{X, X}^* \oplus \mathbb{N}^k \quad \text{"} \langle (f_1)_X, \dots, (f_k)_X \rangle \text{"}$$

This also applies to toric varieties, by taking the pair $(X_\Sigma, X_\Sigma \setminus T)$.

(and also to "toroidal embeddings" i.e.

$U \subseteq X$ that locally look like the embedding of $T_p \subseteq \mathbb{A}^p$).

- log points : $X = \text{Spec } k$ (a point),

$$M = k^* \oplus P \xrightarrow{\alpha} k$$

$$(a, p) \mapsto a \cdot \mathcal{O}_P = \begin{cases} a & \text{if } p=0 \\ 0 & \text{if } p \neq 0 \end{cases}$$

$P = \mathbb{N} \rightsquigarrow$ "standard log point" ("punctured point")



Log structures can be pulled back (and pushed forward):

If $f: X \rightarrow Y$ is a morphism, and (M, d) is a log structure $\xrightarrow{\text{on } Y}$

we can obtain a pullback log str. on X : $M \cong \mathcal{O}_Y$

$f^*M \xrightarrow{f^*d} f^*\mathcal{O}_Y \rightarrow \mathcal{O}_X$ is a pre-log structure, and

we can take the associated log str. $(f^*M)^a \rightarrow \mathcal{O}_X$.

Ex: on $(\mathbb{A}^1, 0)$ we have the divisorial log structure.

The only non-trivial stalk of \mathcal{M} is at 0 , and

$$M_0 \cong \mathcal{O}_{\mathbb{A}^1, 0}^* \oplus \mathbb{N} \leftarrow \langle x \rangle \quad 0 = \{x=0\}$$

We can pull this log str. back to the origin itself along $h: \{0\} \hookrightarrow \mathbb{A}^1$. If you do this, you get exactly the standard log point.

Def: a morphism of log schemes $(X, \mathcal{M}, \alpha) \rightarrow (Y, \mathcal{N}, \beta)$ is a $f: X \rightarrow Y$ + $f^\flat: f^* \mathcal{N} \rightarrow \mathcal{M}$ of log str. on X .

Ex: if $(X, D), (Y, E)$ are SNC pairs, a morphism

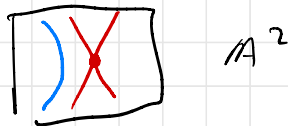
$f: (X, M_D) \rightarrow (Y, M_E)$ is exactly a

$f: X \rightarrow Y$ s.t. $f^{-1}(E) \subseteq D$ ($\Leftrightarrow f(X \setminus D) \subseteq Y \setminus E$)

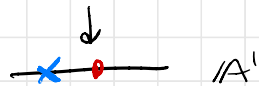
• $\mathbb{A}^2 \rightarrow \mathbb{A}^1$ is of this form

$V(xy=t)$
in \mathbb{A}^3

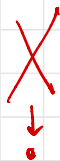
$(x, y) \mapsto xy = t$



"versal deformation of a node"



Can restrict to the origin with the pullback log structures:



the log str. remember "something" about the whole family.

Def: the characteristic sheaf of (M, α) is $\bar{M} = M / \mathcal{O}_X^*$.

This has a "discrete/combinatorial" nature, in the presence of good local models,

Def: a chart for (M, α) is a hom. $P \rightarrow M(X)$ s.t.

$$\underbrace{P \rightarrow M \xrightarrow{\alpha} \mathcal{O}_X}_{\text{pre-log str}} \rightsquigarrow \underbrace{(P)^{\alpha} \xrightarrow{\cong} M \xrightarrow{\alpha} \mathcal{O}_X}_{\text{constant sheaf}}$$

Rule: $P \rightarrow M(X) \rightarrow \mathcal{O}_X(X) \rightsquigarrow k[P] \rightarrow \mathcal{O}_X(X) \rightsquigarrow X \rightarrow \text{Spec } k[P] = \mathbb{A}^1_P$

$P \rightarrow M(X)$ is a chart $\Leftrightarrow X \rightarrow \mathbb{A}^1_P$ is strict, i.e.

the log str on X is the pullback of the divisorial one on \mathbb{A}^1_P .

Def: a log scheme is fs if it locally admits charts from fs monoids.

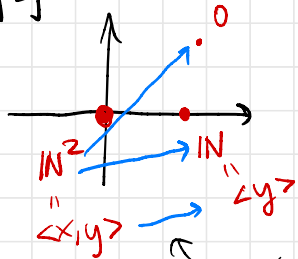
Rule: if $P \rightarrow M(X)$ is a chart, then the stalks of \bar{M} are quotients of P by faces.

If (X, M, α) is fs, then \bar{M} is locally constant on a stratification of X . (this generalizes the orbit

Stratification of a toric variety).

Ex: $\mathbb{A}^2 = \text{Spec} k[\mathbb{N}^2]$

$\bar{M} =$



(analogous picture for \mathbb{A}^p)

kills x which becomes a unit.

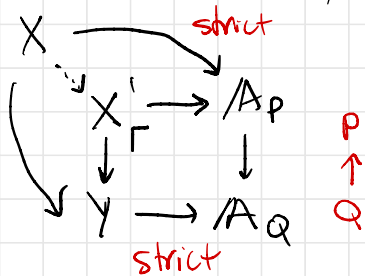
log smoothness and log curves

There is a notion of smoothness for maps of log schemes, which is more flexible than the classical one.

The definition is via an "infinitesimal lifting criterion", but

Thm : $f: (X, M) \rightarrow (Y, N)$ of fs log schemes is log smooth iff locally on X and Y there are charts

$$P \rightarrow M(X), Q \rightarrow N(Y), Q \rightarrow P \text{ s.t.}$$



$X \rightarrow X'$ is (classically)

smooth

(+ technical condition on $Q \rightarrow P$, true e.g.

(if $\text{char} k = 0$ and $Q \leftrightarrow P$)

"log smooth = smooth + combinatorial"

Ex: if $\text{char} k = 0$ and $Q \hookrightarrow P$ is injective $\Rightarrow A_P \rightarrow A_Q$
is log smooth

- $Q=0 \rightsquigarrow A_P \rightarrow \text{Spec} k$ is \checkmark smooth
- in particular $A^2 \rightarrow A^1$ above is log smooth
(it is induced by $\mathbb{N} \rightarrow \mathbb{N}^2$
 $i \mapsto (1, i)$) , and this

stays true for the central fiber



In particular, you can study log smooth curves and their moduli spaces.



Thm: log smooth curves are precisely nodal (marked) curves. Around a node p in a fiber of a family $(C, h_C) \rightarrow (S, h_S)$ of log smooth curves, the family is locally pulled back from $A^2 \rightarrow A^1$ along $(S, h_S) \rightarrow (A^1, 0)$ ($\Leftrightarrow m \in M_S(S)$)

\uparrow
"logarithmic deformation parameter of the node"

The moduli space (stack) of log curves is

isomorphic to the log stack given by the

Deligne-Mumford \rightarrow

DM compactification $\overline{M}_{g,n}$, equipped w/ the
divisorial log str. from the boundary $\overline{M}_{g,n} \setminus M_{g,n}$
(or NC divisor).

Working with log smooth objects ^{sometimes} gives a compact mod. space.

This principle has been applied e.g. to log GW invariants
and log compactification of the universal Jacobian.

Tropical geometry [Chan - lectures on tropical curves]
and their moduli spaces

"piecewise-linear / polyhedral AG"

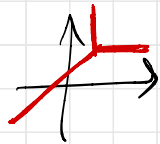
History: if K is a non-archimedean field w/ valuation
 $v: K^* \rightarrow \mathbb{R}$ (e.g. $\overline{K} = \overline{\mathbb{C}((t))} = \bigcup_{n \in \mathbb{N}} \mathbb{C}((t^{1/n}))$)

and a closed $X \subseteq (K^*)^n \rightsquigarrow \text{Trop}(X) = \overline{v(X)} \subseteq \mathbb{R}^n$

$v: (K^*)^n \rightarrow \mathbb{R}^n$
component-wise
valuation.

Thm: $\text{Trop}(X)$ is a (balanced)
polyhedral complex.

e.g. a line $l \in (\mathbb{k}^*)^2$ becomes something like



Curves get tropicalized to metric graphs.

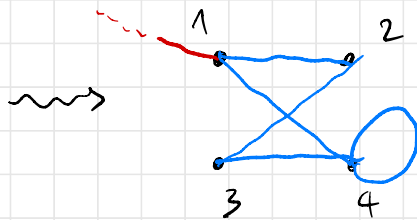
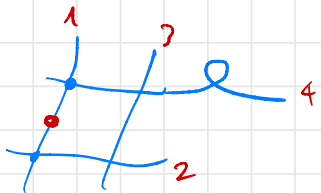
Abstracting:

Def: a tropical curve is a ^{finite} metric graph Γ , possibly with legs (i.e. unbounded edges w/ only 1 endpoint) and a weighting $g: V(\Gamma) \rightarrow \mathbb{N}$.

These things show up as the "dual graph" of a nodal curve:

\mathbb{C} nodal curve $\rightsquigarrow \Gamma(C)$ vertices \leftrightarrow irr. comp. of C
 $g(v) = \text{genus}$ edges \leftrightarrow nodes
 legs \leftrightarrow markings

E.g.:



The length of an edge come from val of a deformation parameter of the corresponding node, because the

Stable curve that you tropicalize will come as the central fiber of a stable model of a smooth curve over \mathbb{K} (whose residue field is the k over which the stable curve lives).

The genus of Γ is $\sum_{v \in V} g(v) + b_1(\Gamma)$. ↙ 1st betti #

There is a moduli space M_g^{trop} of tropical curves of genus g , and this is a ^{generalized} cone complex. There is a tropicalization

map $M_g^{\text{an}} \rightarrow M_g^{\text{trop}}$ ← Berkovich analytification.

There is a version of tropicalization in log geometry, modeled on $X_\Sigma \leftrightarrow \Sigma$ for toric varieties.

Given an fs (X, M, α) , let x_S be the generic point of the log stratum S .

$\overline{M}_{X_S} \rightsquigarrow \sigma_S = (\overline{M}_{X_S})_{\mathbb{R}}^{\vee} = \text{Hom}(\overline{M}_{X_S}, \mathbb{R}_{\geq 0})$ a RPC

if $x_S \rightsquigarrow x_{S'}$ (i.e. $S' \in \overline{S}$), there is a quotient

map $\overline{M}_{X_{S'}} \rightarrow \overline{M}_{X_S}$, that dualizes to

a face inclusion $\sigma_s \hookrightarrow \sigma_{s'}$

e.g.: the quotient $\mathbb{N}^2 \xrightarrow{\pi_2} \mathbb{N}$ on \mathbb{A}^2 that we had before
" $\langle x, y \rangle$ " $\langle y \rangle$

induces to $\mathbb{R}_{\geq 0} \hookrightarrow \mathbb{R}_{\geq 0}^2$
 $\mathbb{R} \mapsto (0, \mathbb{R})$



Define the tropicalization Σ_X as the gluing of all these cones along these face inclusions.

This is a (generalized) cone complex (with integral structure).

Ex: • $\Sigma_{X_{\Sigma}} = |\Sigma|$, but without the embedding in $\mathbb{N}\mathbb{R}$.

• $\Sigma_{\mathbb{A}^p} = P_{\mathbb{R}}^V = \text{Hom}(P, \mathbb{R}_{\geq 0}) = \Sigma_{(\text{Spec } k, k^* \otimes P)}$

Tropicalizing log curves gives (families) of tropical curves. When you take the dual graph, the "length" of an edge is the log. def. parameter of the corresponding node.

$$\underline{Ex}: \mathbb{A}^2 \rightarrow \mathbb{A}^1$$

$$(x, y) \mapsto xy$$

$$xy = t$$

\rightsquigarrow

$$\Sigma_{\mathbb{A}^2} = (\mathbb{R}_{>0})^2$$

\downarrow

$$\Sigma_{\mathbb{A}^1} = \mathbb{R}_{>0}$$



This stuff was fundamental for various recent developments e.g. in log GW theory, the Gross-Siebert program, and compactifications of the universal Jacobian.